

# Fixed Point Analysis of Single Cell IEEE 802.11e WLANs: Uniqueness, Multistability and Throughput Differentiation

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# Outline of Talk

- 1 Introduction
- 2 Fixed Point Based Saturation Analysis
- 3 Homogeneous Nodes: Uniqueness, Multistability, Unfairness
- 4 Nonhomogeneous Nodes

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- We will focus on the Distributed Coordination Function (DCF) algorithm of IEEE 802.11: a form of CSMA/CA
- Why study such a MAC.... it is an old and well established technique?
- Many current and emerging systems utilise this MAC
- In addition to the IEEE 802.11 series of networks
  - Sensor networks: SMAC, and IEEE 802.15.4
  - Availability of inexpensive chip-sets is driving new applications
    - e.g., rural access networks based on long distance IEEE 802.11 links
- A variety of new questions are being asked
  - Service differentiation, delay analysis, admission control, power control
  - Enterprise WLAN planning; automatic channel allocation
  - Enterprise WLAN performance monitoring and management
    - How should "packet sniffer" measurements be used?
  - Multihop IEEE 802.11 networks

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# Traffic Models Used in Our Analyses

- Nodes saturated: i.e., “infinite” backlog of packets
  - Saturation analysis is an important first step
  - Provides a measure of network capacity
    - Sometimes can be proven to be a sufficient stability condition
  - Directly useful in some modeling situations
  - By-products of saturation analysis have been found to work extremely well in unsaturated models (e.g., modeling with VoIP and TCP)
- Our work has appeared in: IEEE Infocom 2005, IEEE Sigmetrics 2005 IEEE/ACM Tr. on Networking 2007, and IEEE/ACM Tr. on Networking 2008
  - The present talk will summarise some of this work
- The following have been studied, but will not be reported in this talk
  - TCP controlled file transfers
  - Packet voice telephony (CBR packet streams)
  - Streaming video
  - Combinations of these traffic types

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# Homogeneous vs. Nonhomogeneous

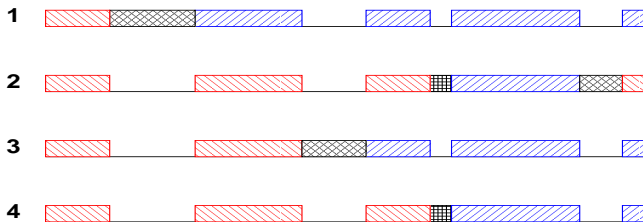
- **Homogeneous:** All nodes have the same access parameters
  - This is the situation in IEEE 802.11b
- **Nonhomogenous:** Nodes can have difference access parameters
  - This would be done to provide service differentiation
  - The IEEE 802.11e standard allows this
- We will begin with the analysis of the backoff process in the homogeneous case

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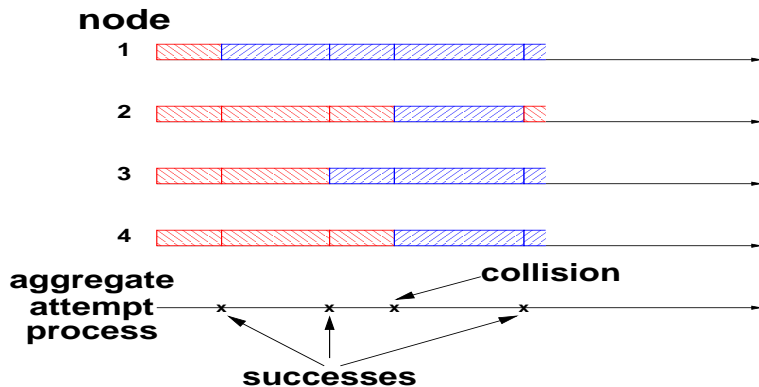
# Saturated Nodes: Backoffs, Transmissions, Collisions

node



- Channel alternates between activity and back-off periods
- Backoff time is slotted and freezes during channel activity
- Activity period depends on which nodes cause the activity

# Focusing Only on the Backoff Durations



- This evolution depends only on the back-off parameters, and the collision and success events
- The nodes' backoff evolutions are coupled due to collisions
- Can reconstruct original process from this process

# A Detailed Markov Chain

- There are  $n$  nodes, with saturated queues
- Time evolves over system slots,  $t \in \{0, 1, 2, \dots\}$
- Let  $Z_i(t)$  be the backoff stage of node  $i$  in slot  $t$ 
  - $Z_i(t) \in \{0, 1, \dots, K\}$
- Let  $Y_i(t)$  be the residual backoff of node  $i$  in slot  $t$
- Then  $\{((Z_1(t), Y_1(t)), (Z_2(t), Y_2(t)), \dots, (Z_n(t), Y_n(t))), t \geq 0\}$  is a Markov chain
- Too complex to analyse, even if we take the backoff durations to be geometrically distributed

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# Attempt and Collision Rates

- Let  $\beta$  be a node's attempt probability (in a backoff slot)
  - i.e.,  $\beta$  is the attempt rate of a node, per slot, averaged over all back-off slots
  - Let  $A_i(t)$  = number of attempts by Node  $i$  in  $[0, t]$ , then

$$\beta = \lim_{t \rightarrow \infty} \frac{A_i(t)}{t}$$

- Let  $\gamma$  be the collision probability seen by a node
  - i.e.,  $\gamma$  is the fraction of a node's attempts that collide with other attempts
  - Let  $C_i(t)$  = number of collisions seen by Node  $i$  in  $[0, t]$ , then

$$\gamma = \lim_{t \rightarrow \infty} \frac{C_i(t)}{A_i(t)}$$

- Thus, these are time average rates
- Two equations can be written relating these quantities
  - Bianchi, IEEE JSAC'00; our generalisation, IEEE Infocom'05, IEEE/ACM TON 2007

# A Fixed Point Equation

- Isolate a single node and view the remaining  $n - 1$  as an “environment”
- “Decoupling” approximation
  - From the point of view of a node, the other nodes are viewed as attempting independently in each slot with probability  $\beta$
- The node “response” equation
  - $\beta = G(\gamma)$ : models how a node adjusts to collisions
- The coupling equation
  - $\gamma = \Gamma(\beta) := 1 - (1 - \beta)^{n-1}$
  - Consequence of the “decoupling” approximation
- This yields the following, natural, fixed point equation

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# The Backoff Parameters for Each Node

- Homogeneous case
- The following are the back-off parameters for each packet attempted by the MAC at a node
- $K$  := At the  $(K+1)$ th attempt either the pkt succeeds or is discarded
  - $K$  is the maximum number of retries
- $b_k$  := The mean back-off at the  $k$ th attempt of a pkt,  $0 \leq k \leq K$

# The Node Response Formula $G(\gamma)$

$$G(\gamma) = \frac{1 + \gamma + \gamma^2 \cdots + \gamma^K}{b_0 + \gamma b_1 + \gamma^2 b_2 + \cdots + \gamma^k b_k + \cdots + \gamma^K b_K}$$

- The formula captures all backoff parameters
- The distribution of the back-off durations does not matter
- No need to keep track of the residual backoff
  - ... as done by Bianchi (and several other researchers)

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# Uniqueness Theorem 1

- We have the fixed point equation  $\gamma = \Gamma(G(\gamma))$

## Lemma

*$G(\gamma)$  is nonincreasing in  $\gamma$  if  $b_k, k \geq 0$ , is a nondecreasing sequence.*

- $\Gamma(\beta)$  is increasing in  $\beta$

## Theorem

*$\Gamma(G(\gamma)) : [0, 1] \rightarrow [0, 1]$ , has a unique fixed point if  $b_k, k \geq 0$ , is a nondecreasing sequence.*

- In practice, in IEEE 802.11,  $b_k$  is a nondecreasing sequence
- The fixed point is then taken as the operating point of the system

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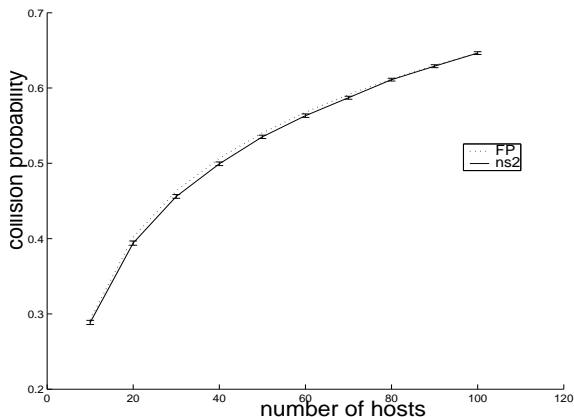
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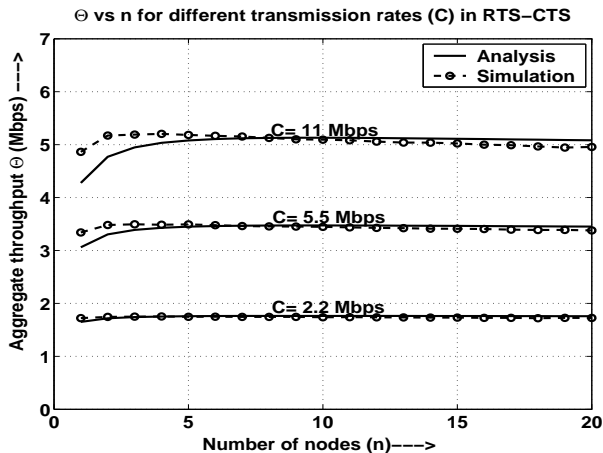
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# Comparison with Simulation: $\gamma$ vs. $n$



- IEEE 802.11 parameters used:  $b_0 = 16$  slots,  $b_k = 2b_{k-1}, 1 \leq k \leq 7$
- The fixed point analysis provides a very good approximation for a wide range of values of the number of nodes

# Comparison with ns2 Simulations



- Note how the saturation throughput is fairly constant with the number of nodes
- Motivates a state dependent service rate model

# But Will This Analysis Always Work?

- Is it accidental that it works for the IEEE 802.11 parameters?
- Why does it work?
- Can it be applied to other 802.11-like MACs?
- Next, we provide some partial answers...

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# Homogeneous Nodes: Unbalanced Fixed Points

- A fixed point  $(\gamma_1, \gamma_2, \dots, \gamma_n)$  is called **balanced** if all the coordinates are equal
- A balanced fixed point would be expected to capture the steady state performance
- But can there be unbalanced fixed points?
  - i.e., a vector  $(\gamma_1, \gamma_2, \dots, \gamma_n)$ , s.t., for  $1 \leq i \leq n$ ,

$$\gamma_i = 1 - \prod_{j \neq i} (1 - G(\gamma_j))$$

with all coordinates not equal

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- Would suggest multistability or short term unfairness

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# Example 1: Homogeneous System, Unbalanced Fixed Points

- There are  $n = 10$  nodes

$$G(\gamma) = \frac{1 + \gamma + \gamma^2 + \gamma^3 + \dots}{1 + \gamma + \gamma^2 + \gamma^3 + 64(\gamma^4 + \gamma^5 + \dots)}$$

- i.e.,  $K = \infty$ ,  $b_0 = b_1 = b_2 = b_3 = 1$ ,  $b_4 = b_5 \dots = 64$ 
  - $b_0, b_1, b_2, b_3$  arise from a point mass at 1
- Intuitively, one node will go on succeeding while all others will be in a long back-off
  - The favoured node will keep changing randomly

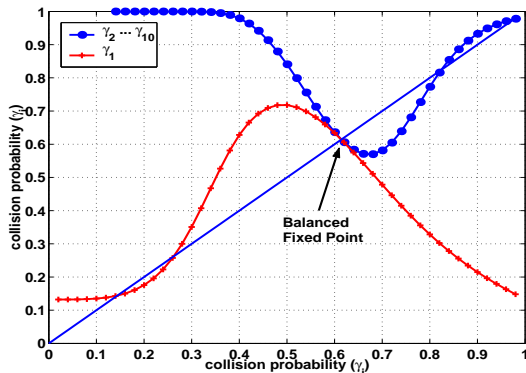
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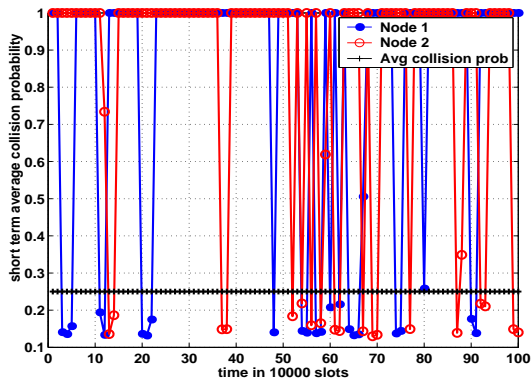
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## (... contd.) Example 1: Unbalanced FPs

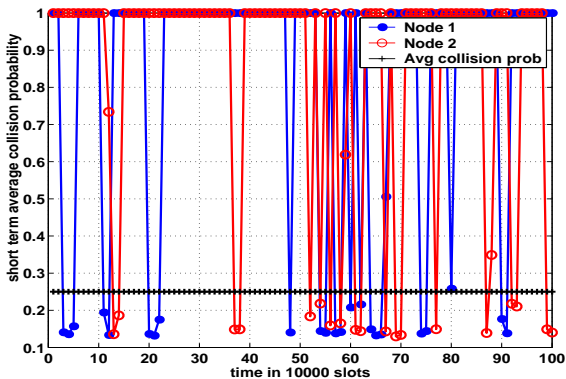


- The balanced fixed point  $\gamma_i = 0.62, 1 \leq i \leq 10$ , is shown
  - Recall, from earlier theory, that a balanced fixed point must exist
- An unbalanced FP is  $\gamma_1 = 0.14, \gamma_i = 0.97, 2 \leq i \leq 10$

## (... contd.) Example 1: Multistability



- Short term average collision rate, over 10,000 slot segments
- The node that is being favoured gets a small collision rate
  - while the rest see a collision rate close to 1



- Time average collision rate is 0.25
  - which **does not** match the balanced fixed point: 0.62

## Example 2

$$G(\gamma) = \frac{1 + \gamma + \gamma^2 + \dots + \gamma^7}{1 + 3\gamma + 9\gamma^2 + 27\gamma^3 + \dots + 2187\gamma^7}$$

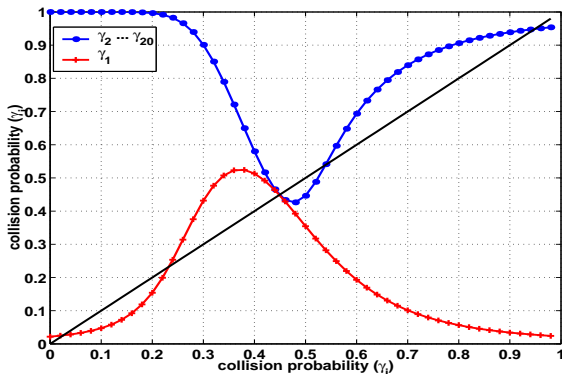
- $n = 20$  homogeneous nodes, with the above form of  $G(\cdot)$
- This is in the IEEE 802.11 DCF back-off expansion framework
- ... with  $K = 7$ ,  $b_0 = 1$ ,  $p = 3$  and  $b_k = p^k b_0$  for all  $0 \leq k \leq K$
- Again there are multiple nonhomogeneous fixed points
- ... and the balanced fixed point does not capture the steady state performance

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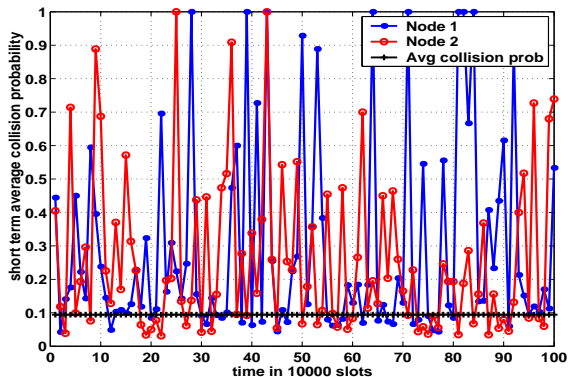
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## (... contd.) Example 2: Unbalanced FPs



- The balanced fixed point  $\gamma_i = 0.45, 1 \leq i \leq 20$ , is shown
- Unbalanced FPs are also evident

## (... contd.) Example 2: Multistability



- Short term average collision rate, over 10,000 slot segments
- The unfairness is less than in Example 1

$$G(\gamma) = \frac{1 + \gamma + \gamma^2 + \dots + \gamma^K}{b_0(1 + p\gamma + p^2\gamma^2 + \dots + p^K\gamma^K)}$$

## Theorem

*For the function  $G(\cdot)$ , defined above, if  $K \geq 1$ ,  $p \geq 2$  and  $b_0 > 2p + 1$ , then the system  $\underline{\gamma} = \mathbf{\Gamma}(\mathbf{G}(\underline{\gamma}))$  has a unique fixed point. This fixed point is (of course) balanced.*

- Example 2 did not satisfy this sufficient condition
  - ... and we had found that it had unbalanced fixed points

# The IEEE 802.11 Standard DCF

- The above result easily extends to the IEEE 802.11 standard where  $b_k = p^k b_0$  for  $0 \leq k \leq m \leq K$  and  $b_k = p^m b_0$  for  $m < k \leq K$
- Hence for the IEEE 802.11 standard there is a unique balanced fixed point and there can be no unbalanced fixed points
- The fixed-point analysis has been well known to estimate the saturation performance very well

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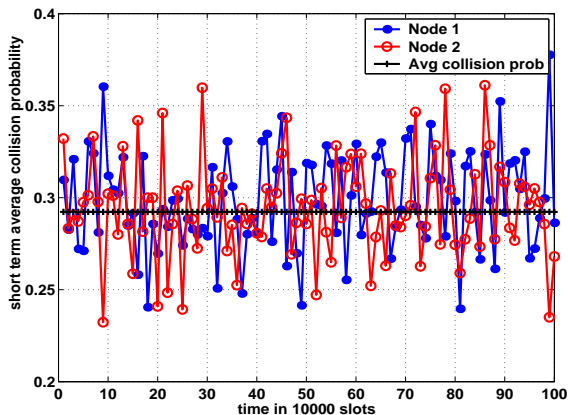
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- The fixed-point analysis has been well known to estimate the saturation performance very well

## Example 3

$$G(\gamma) = \frac{1 + \gamma + \gamma^2 + \dots + \gamma^7}{16 + 32\gamma + 64\gamma^2 + \dots + 2048\gamma^7}$$

- $n = 10$  homogeneous nodes with the above  $G(\cdot)$
- Satisfies the condition of the above theorem
  - The parameters are close to those used in IEEE 802.11
- Hence unique fixed point; this fixed point is balanced

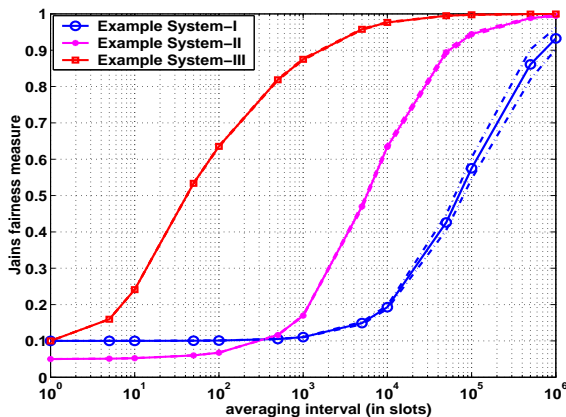
## Example 3: Lack of Short Term Unfairness



- Also, the balanced fixed point captures the steady state performance well ( $\gamma = 0.29$  from fixed point and from simulation)

# Fairness Comparison

- Fairness measure:  $\frac{1}{n} \frac{(\sum_{i=1}^n \tau_i)^2}{\sum_{i=1}^n \tau_i^2}$ , where  $\tau_i$  is the average throughput of node  $i$  over the averaging interval



# Some Lessons Learnt

- A network with homogeneous nodes can have unbalanced fixed points
- Nonunique fixed points appear to imply multistability
  - ... which manifests itself as unfairness over rather long periods
- When there are nonunique fixed points the balanced fixed point does not capture steady state performance
  - ... which suggests that fixed point analysis will be accurate only when there is a unique, balanced fixed point

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# Uniqueness Theorem 3

- $n$  contending nodes, each with one queue whose  $G_i(\cdot)$  function is of the form

$$G_i(\gamma_i) = \frac{1 + \gamma_i + \gamma_i^2 + \cdots + \gamma_i^K}{b_{0_i}(1 + p_i\gamma_i + p_i^2\gamma_i^2 + \cdots + p_i^K\gamma_i^K)}$$

- IEEE 802.11e standard for differentiated services with DCF

## Theorem

*If  $G_i(\cdot)$  satisfies  $K_i \geq 1$ ,  $p_i \geq 2$  and  $b_{0_i} > 2p_i + 1$ , there exists a unique fixed point for the system of equations  $\gamma_i = 1 - \prod_{j \neq i} (1 - G_j(\gamma_j))$  for  $1 \leq i \leq n$ .*

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  - ... a low priority queue can defer access in favour of higher priority queues
  - ... by waiting longer than DIFS after each channel activity period
- A decoupling approximation leads to a vector fixed point equation

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# Comments on the Fixed Point Approach

- Easy to develop the equations and obtain numerical results
  - ... but needs to be used with care, as we have seen
- A formal linkage between the system and the fixed point equations is now emerging
  - Bordenave, McDonald, Proutiere (Allerton 2005): Prove that the back-off stage process decouples in a certain limiting regime
- Such fixed point techniques exist for other problems
  - F.P. Kelly: alternate routing in circuit switched networks
  - A. Stolyar: analysis of input queued packet switches

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# Comments on the Fixed Point Approach

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